ECS165A

Milestone 2 Overview

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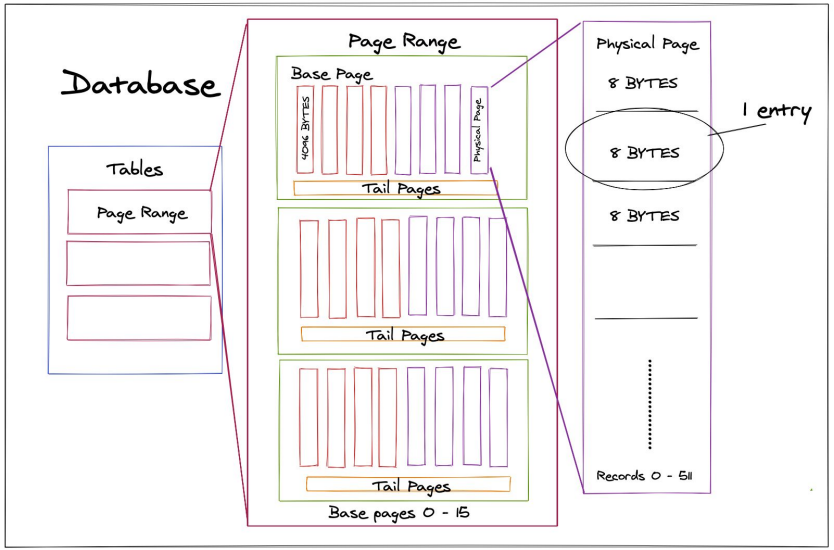
Objectives

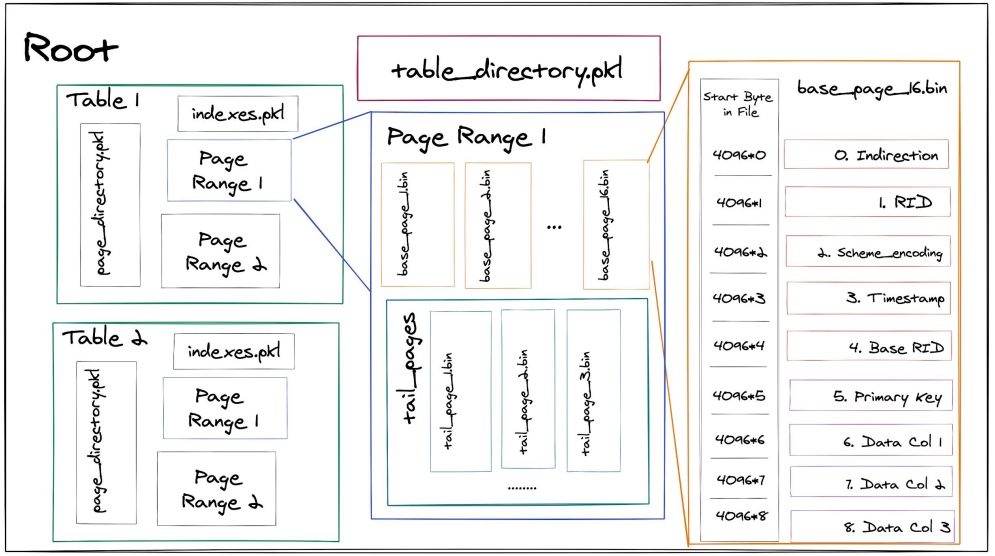
- Disk Structure
- Bufferpool Design
- Eviction Policy
- Indexing
- Merge
- Performance
- Questions
- Demo

The Path to Durability

Milestone 1



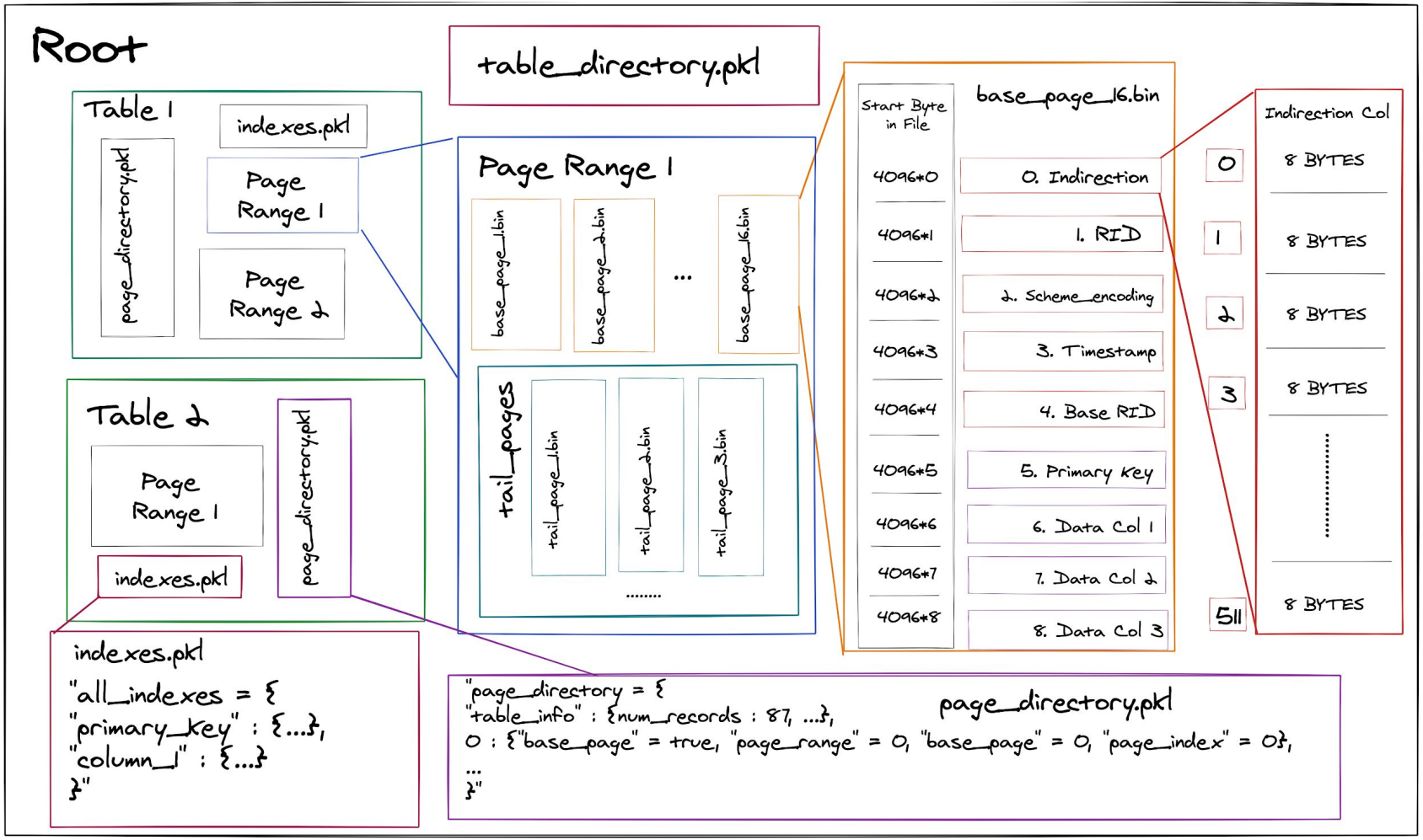
Milestone 2



Volatile Memory

Non-Volatile Memory

Disk Structure



Bufferpool Design

Bufferpool

*A frame loads 1 base or tail page

```
frame_directory = {
(table_name, page_range, base/tail page index, is_base_record) : frame_index
}
```

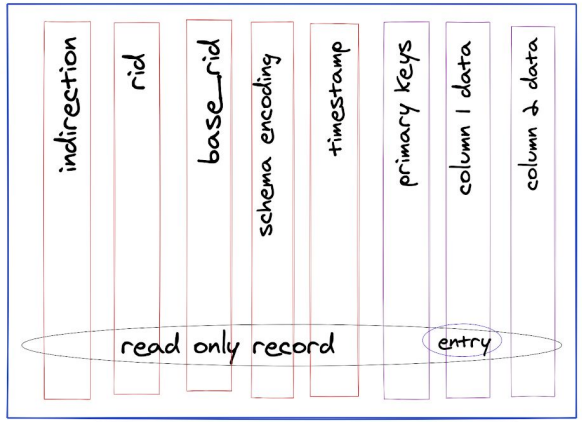
```
Main Methods
evict_page()
load_page()
commit_page()
at_capacity()
is_record_in_pool()
add_frame_to_directory()
commit_all_frames()
```

Frame Array

Frame 1	Frame 2	Frame 3	Frame 4	...

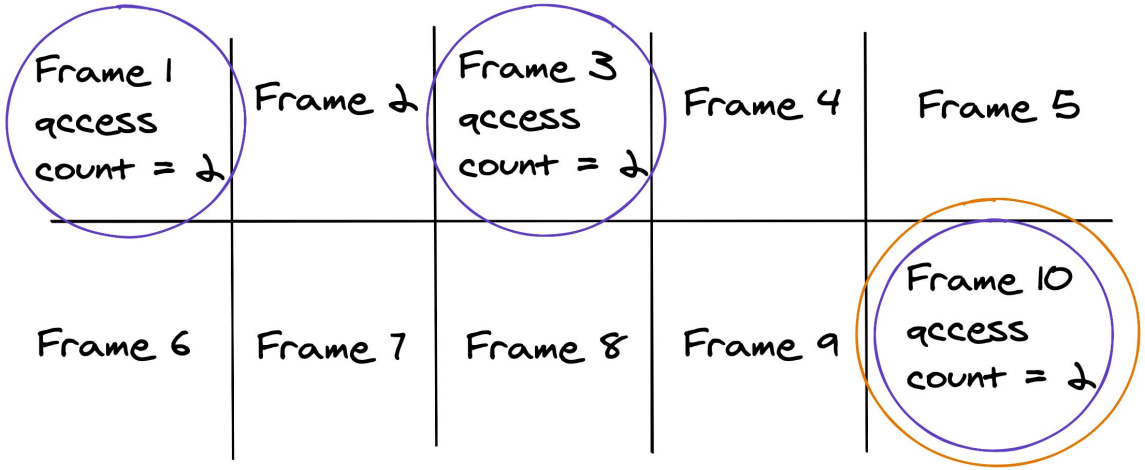
*Bufferpool size is set in the configuration file

Frame
- all_columns
- pin
- dirty_bit
- time_in_bufferpool
- access_count
- path_to_page_on_disk
- frame_key



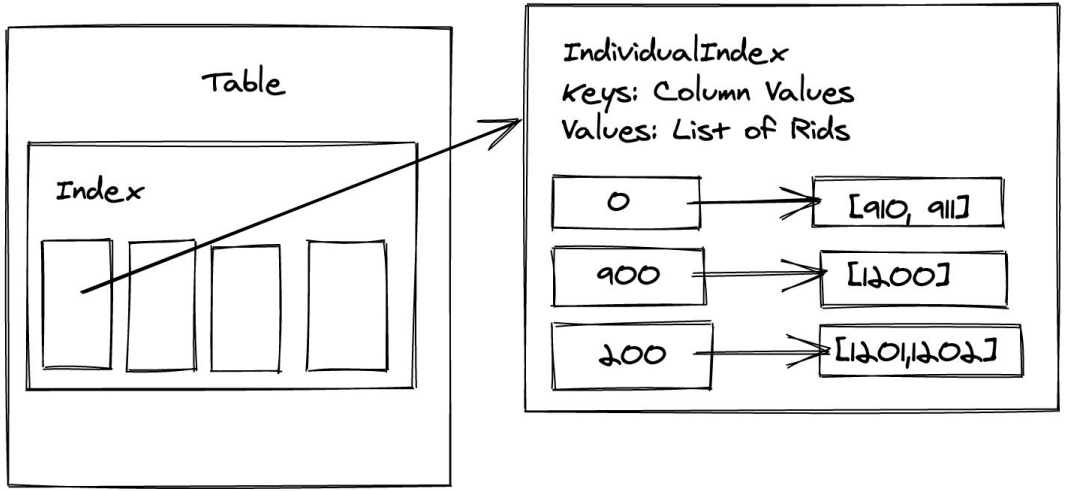
Eviction Policy

- Synthesis of LRU and LFU policies
  - Among the least frequently used records, evict the least recently used
- Chose this method primarily for speed & simplicity
- Does not distinguish between privileged and unprivileged data

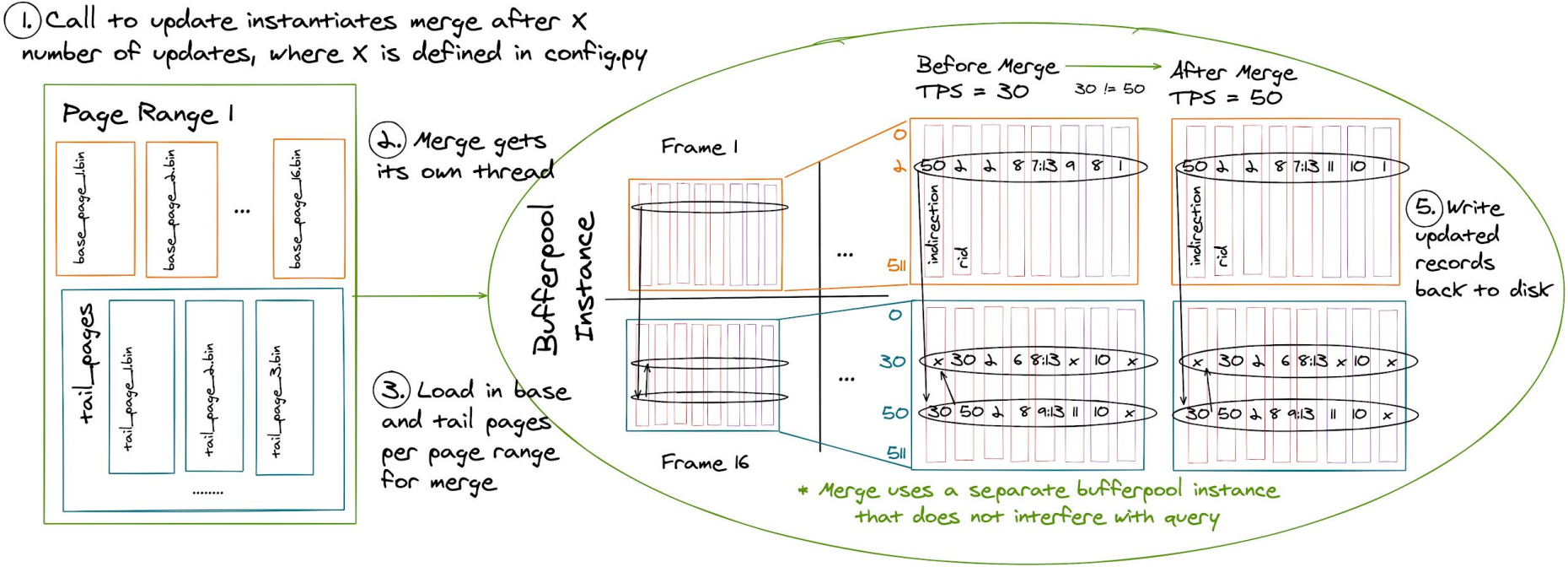


Indexing

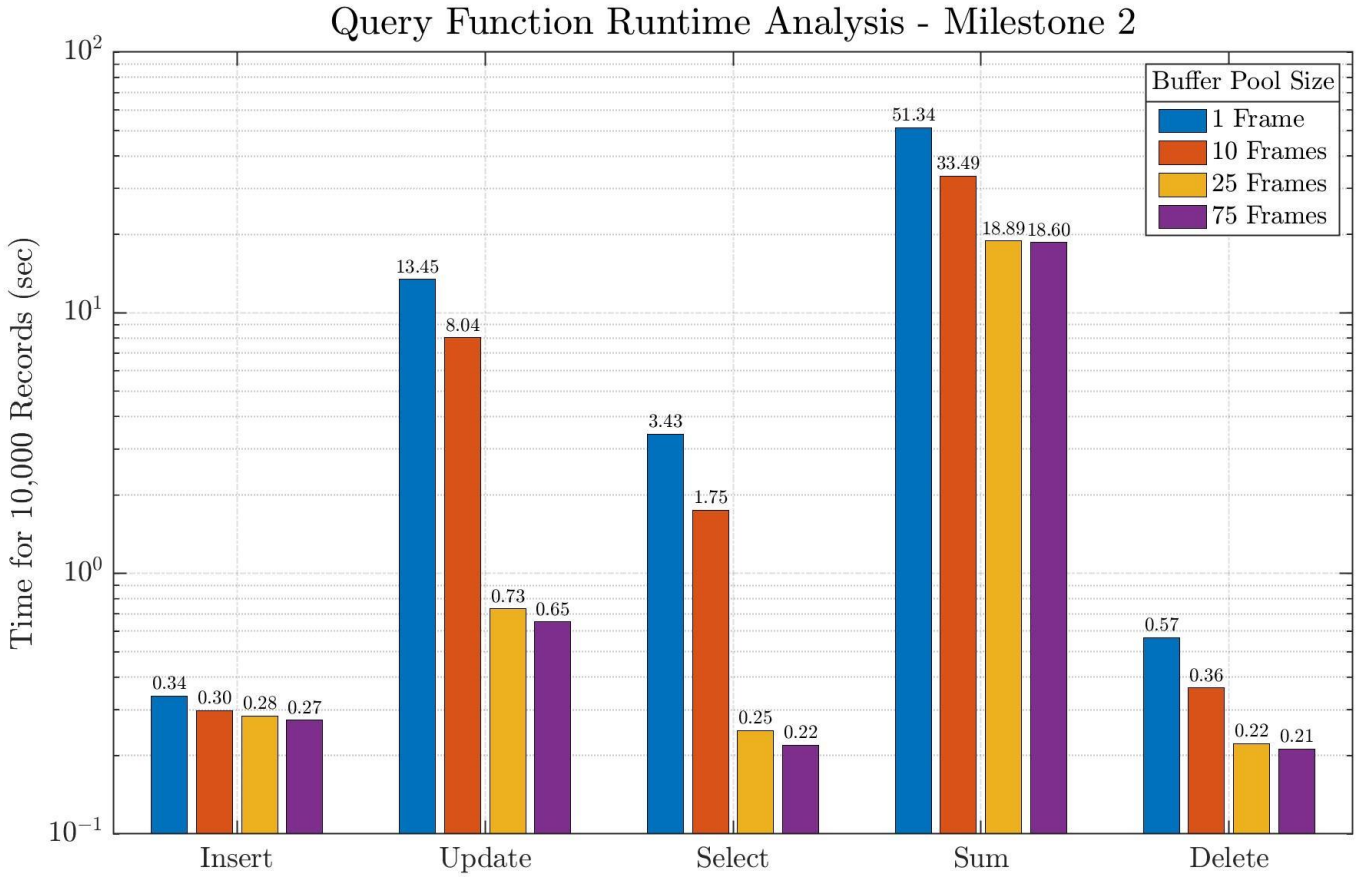
- Hash index
- Maps column values to list of RIDs
- Used for select and update
- Index created automatically for primary key
- Bufferpool frames are committed before index creation
- Index persisted as .pkl file



Merge - General Flow



Optimizations & Performance



*These times are based on 10 run averages using the provided __main__.py

Specs: 6 core Intel Core i7, 2.6 GHz, 256KB l2 cache per core, 12 MB L3 Cache, 16GB Memory

Questions

Demo