



# L-Store Milestone 2

**ECS 165A**

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# Durability & Bufferpool Extension

Save & Load Pages from Persistent Storage

Pinning & Unpinning Pages

# Data Reorganization

Merge

# Indexing

B-Tree Index

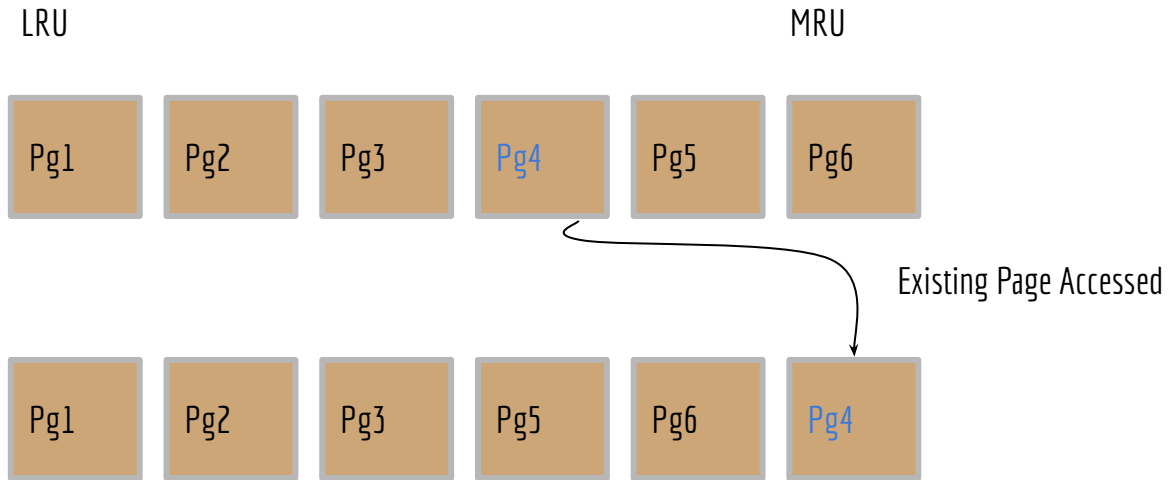


# Durability & Bufferpool Extension



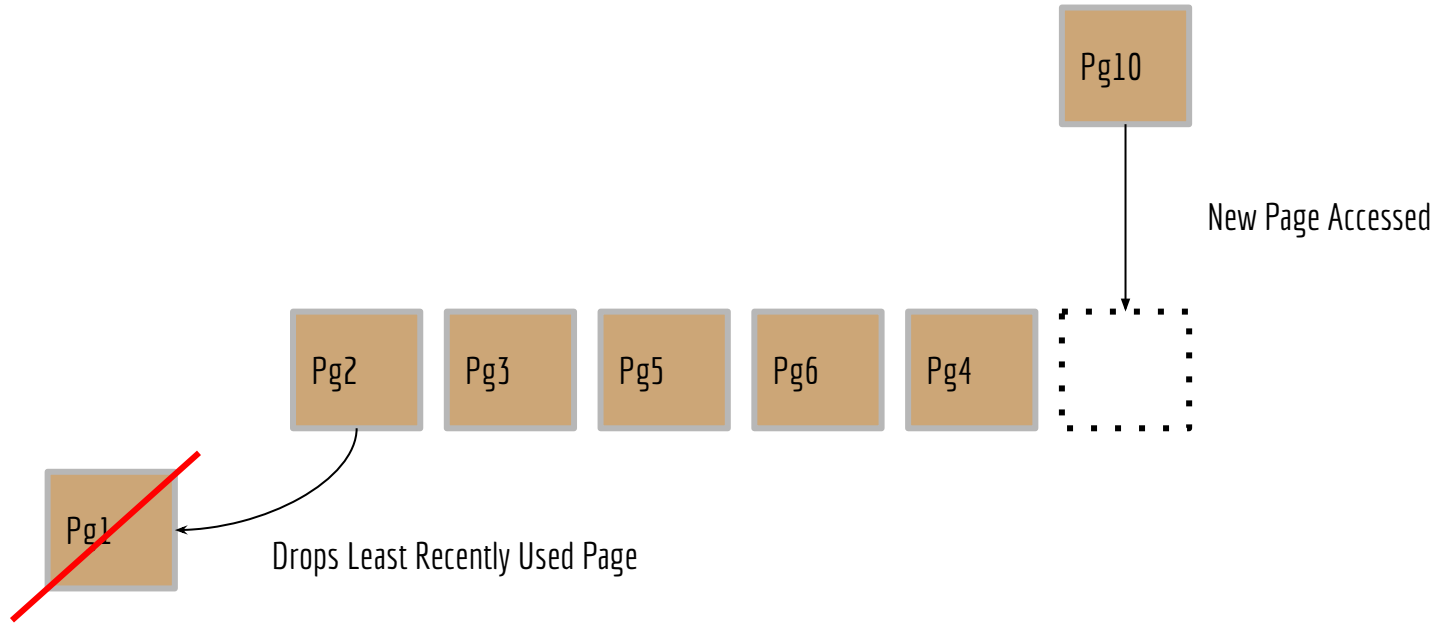
# Persistent Storage - Accessing within Bufferpool

Ordered Dictionary



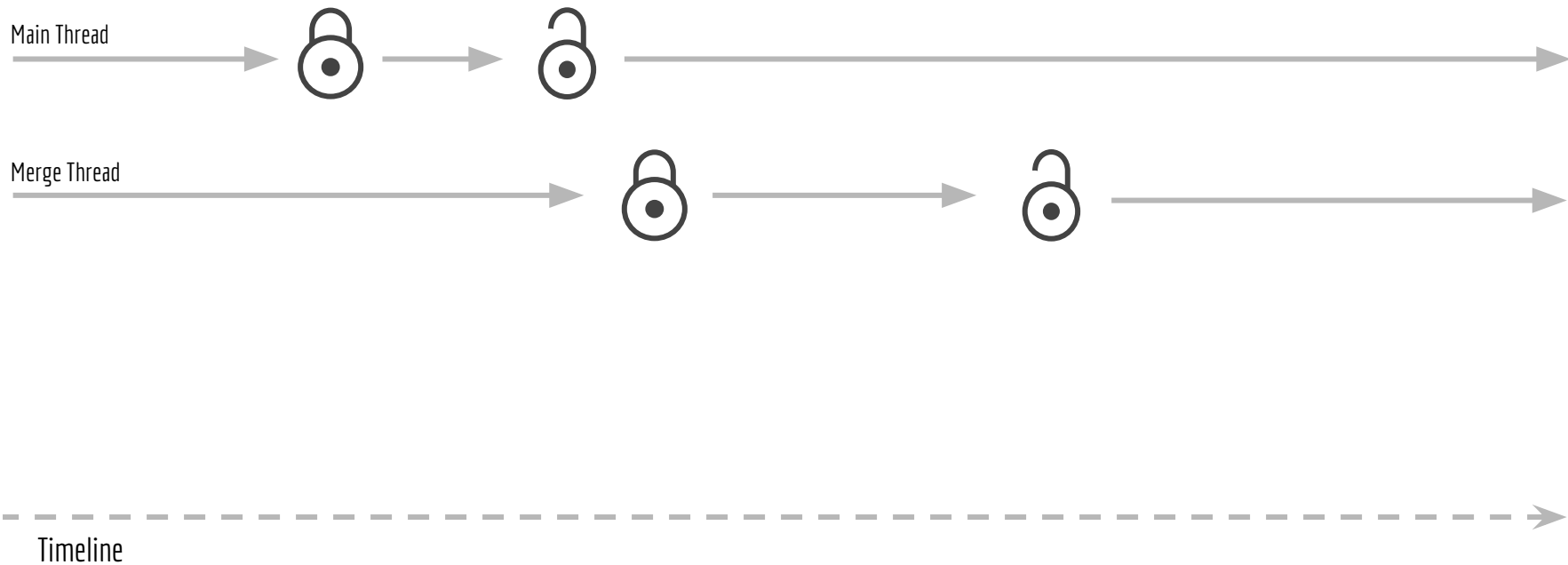
# Persistent Storage - Accessing at Capacity

Ordered Dictionary

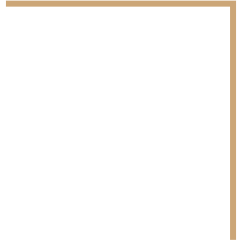
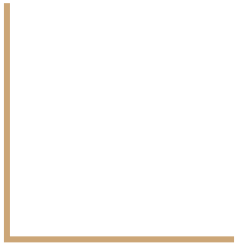


# Pinning & Unpinning

Utilizing Binary Semaphores



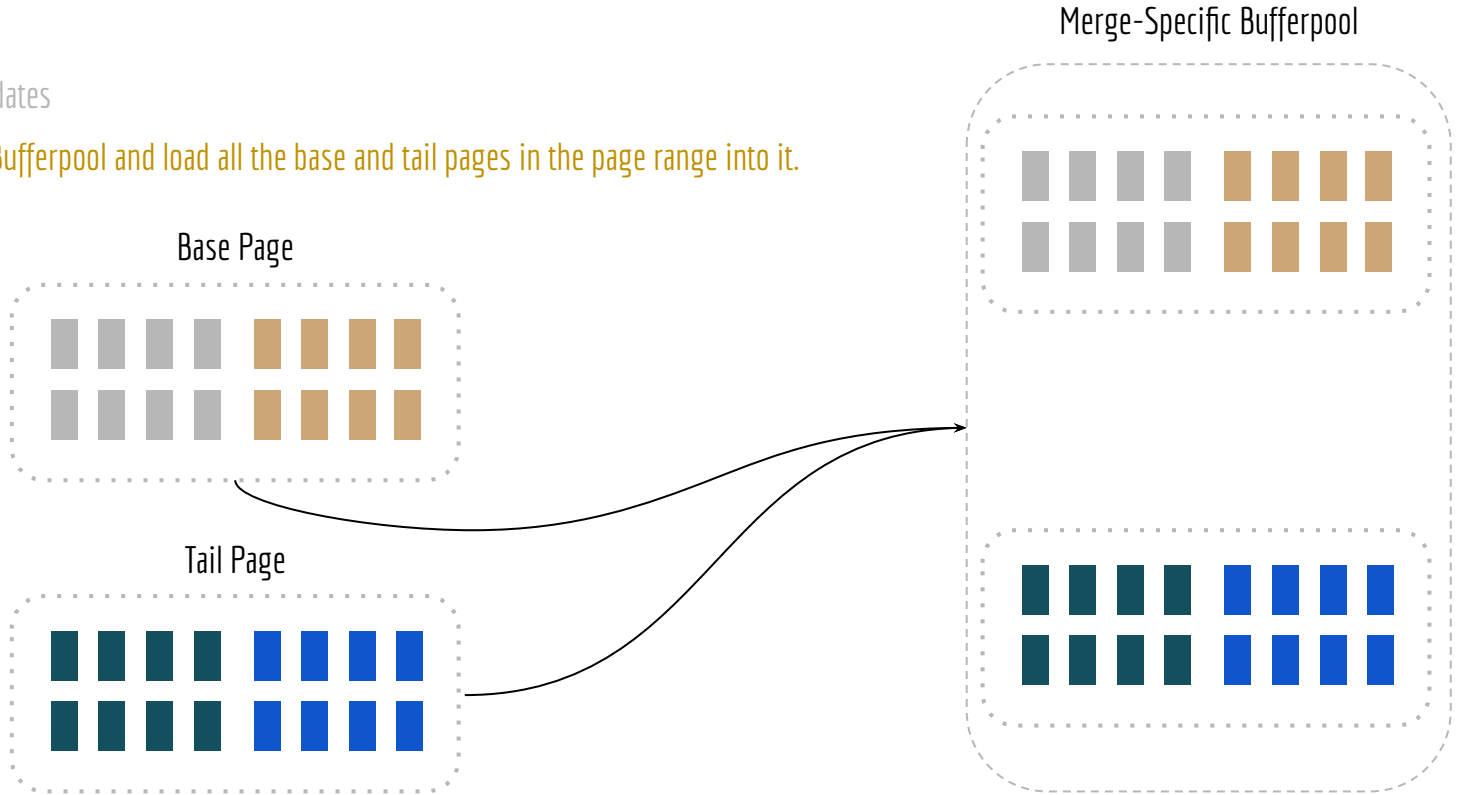
# Data Reorganization



# Merge

Default Threshold: 2048 updates

1. Create a Merge-Specific Bufferpool and load all the base and tail pages in the page range into it.

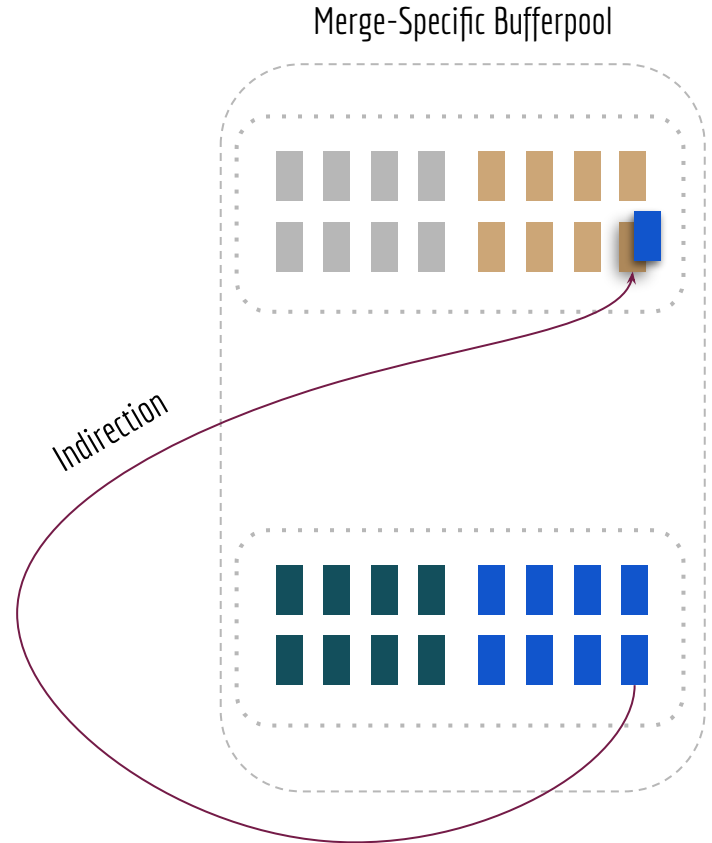
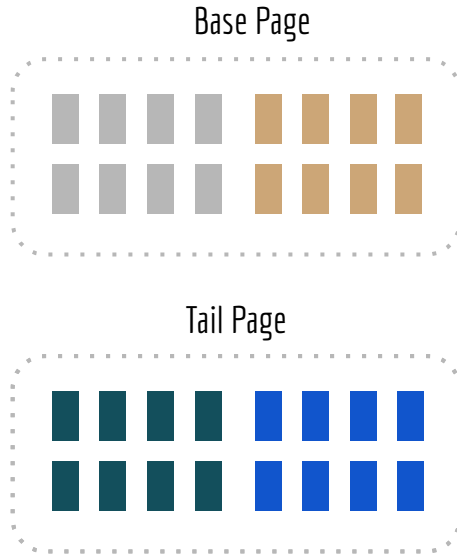




# Merge

Default Threshold: 2048 updates

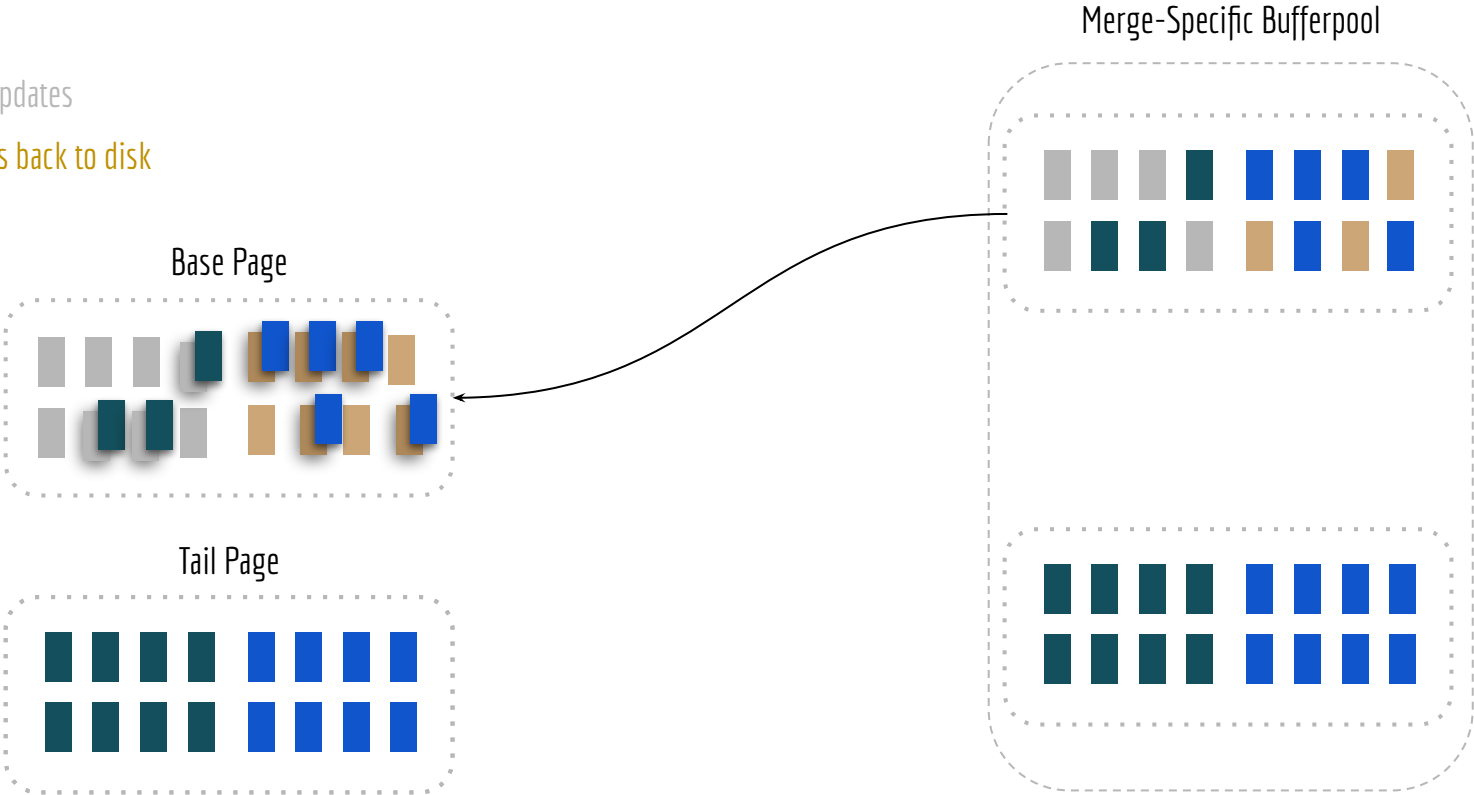
2. Consolidate all the records from bottom up



# Merge

Default Threshold: 2048 updates

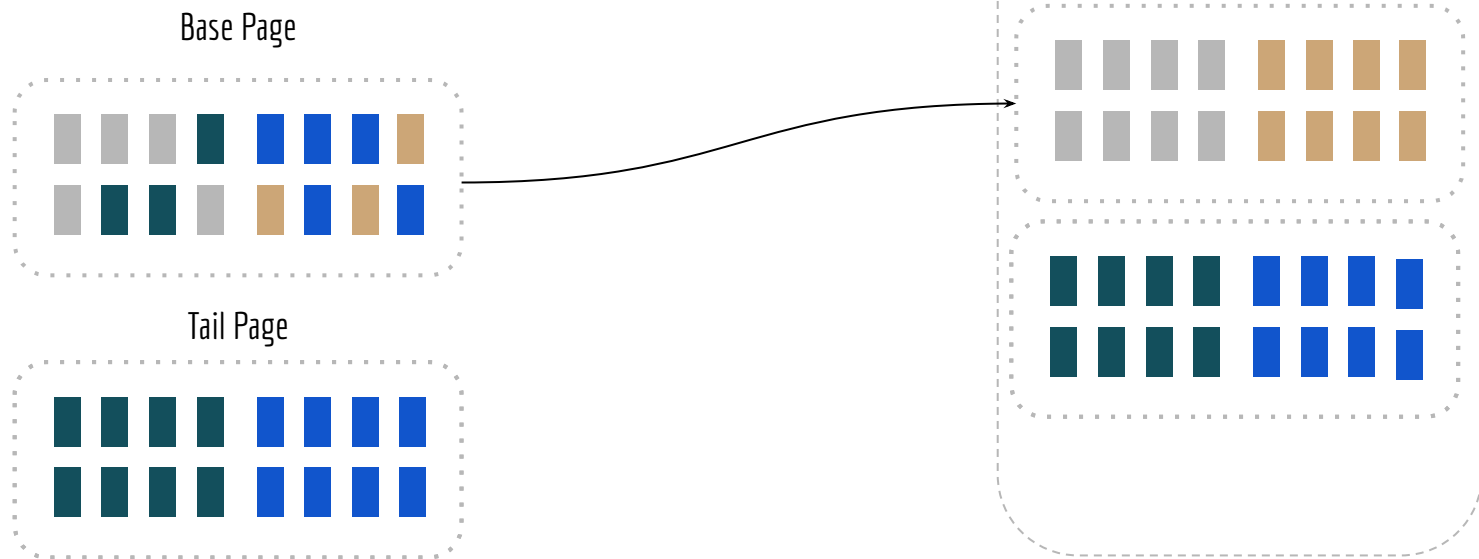
3. Write the updated pages back to disk



# Merge

Default Threshold: 2048 updates

4, Load newly updated pages from disk into main bufferpool to refresh it



# Indexing



# Indexing

Our BTree previously utilized OOBTree, where both the keys & values are Objects. We switched it to IOBTree, where the key is an Int and value is an Object.

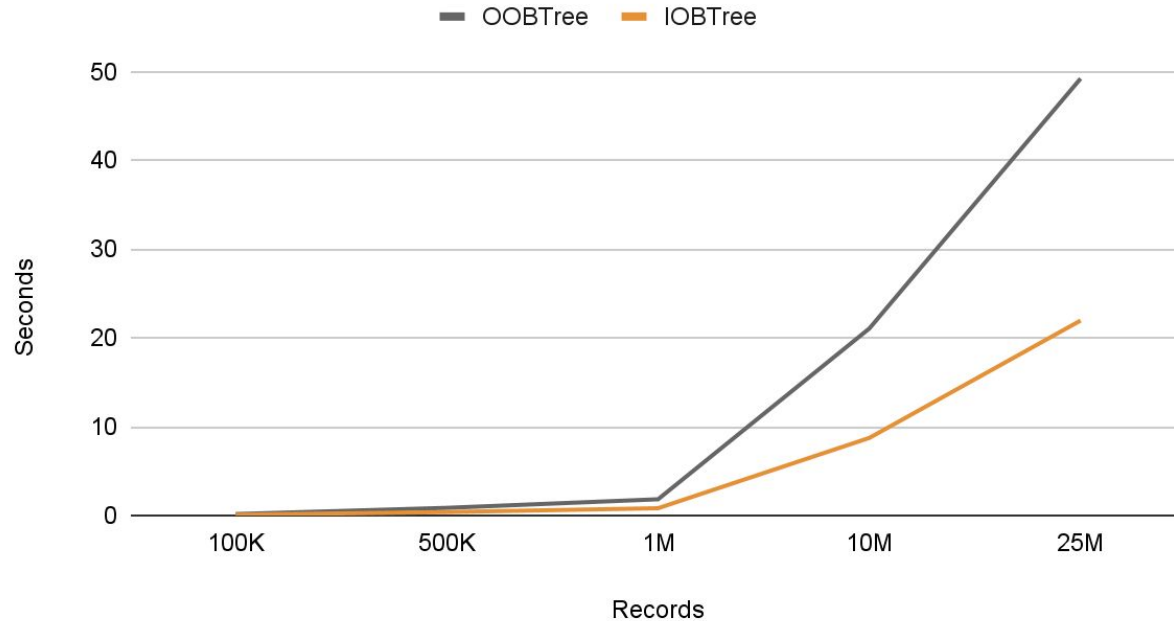
IOBTree is superior to OOBTree both in terms of **memory usage** and **time**.

## Benchmark Program:

1. Insert # key-value pairs
2. Select # key-value pairs
3. Delete # key-value pairs

# OOBTree vs IOBTree

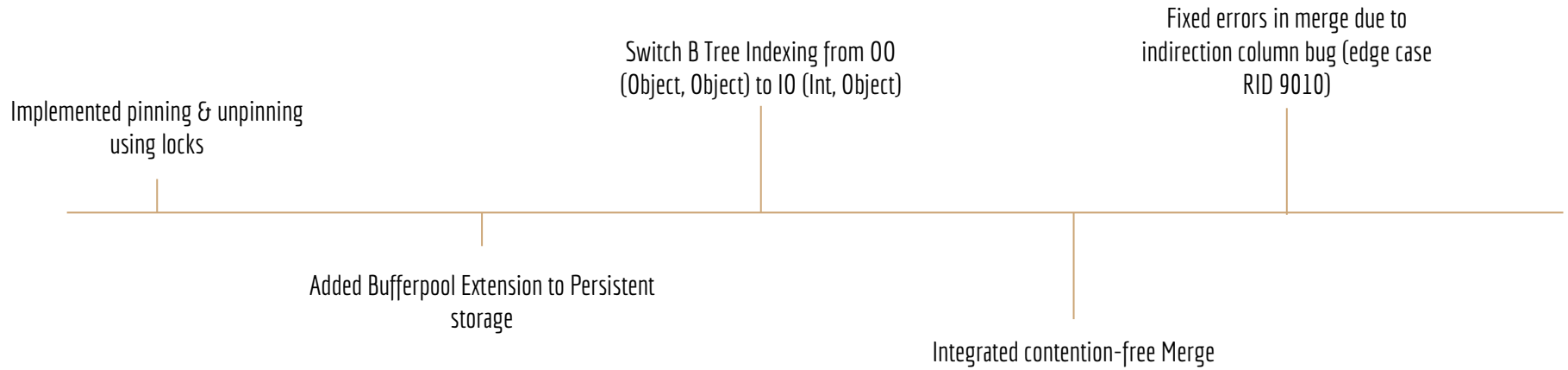
## Time Comparison



# Indexing

Additional implementation

# Roadmap





# Improvements for Milestone 3

- Complete support for multithreading
- Organize code
- Test different values in configurations for better performance

Thank you!

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