



The Journey of Building Global-Scale ResilientDB Blockchain Fabric

Mohammad Sadoghi

URCS Seminar Series University of Rochester November 5, 2021



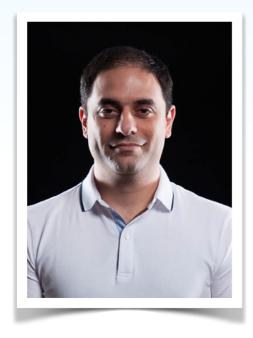
Mohammad Sadoghi

Exploratory Systems Lab Department of Computer Science









Mohammad Sadoghi (Principal Investigator)



Thamir Qadah, PhD (Distributed & Coordination-free Concurrency)



Haojun (Howard) Zhu, MSc (Re-Configurable Consensus Protocols)



Alejandro Armas, BSc (Re-engineering ResilientDB Toolkits)



Shubham Pandey, MSc (Scaling Fabric via RDMA)



Jelle Hellings, PostDoc Suyash Gupta, PhD (Fault-tolerant Complexity Analysis) (Scalable Consensus Meta-Protocols)



Sajjad Rahnama, PhD (Global Scale Consensus)



Dakai Kang, BSc (View-change-less Protocols)

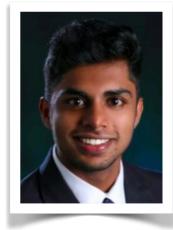


Rohan Sogani, MSc

(Scaling Fabric via Sharding)



Xinyuan Sun, MSc (Scaling Fabric via RDMA)



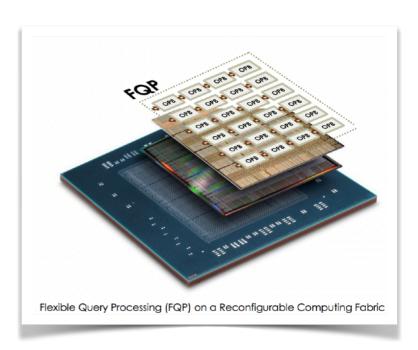
(Scaling Fabric via Sharding)

Dhruv Krishnan, MSc

Priya Holani, MSc

(Scaling Fabric via Sharding)

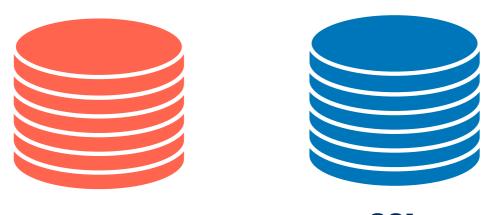




FPGA Acceleration: FQP (Flexible Query Processor) [VLDB'10, ICDE'12, VLDB'13, ICDE'15, SIGMOD Record'15, ICDE'16, USENIX ATC'16, ICDCS'17, ICDE'18, TKDE'19]

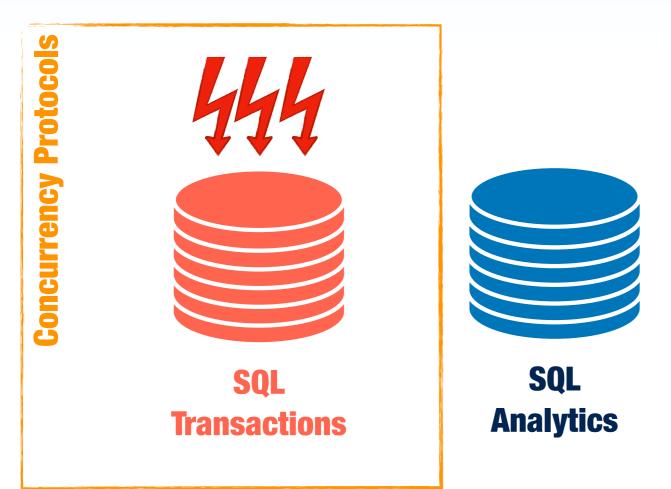
SQL

Analytics



SQL Transactions **SQL** Analytics

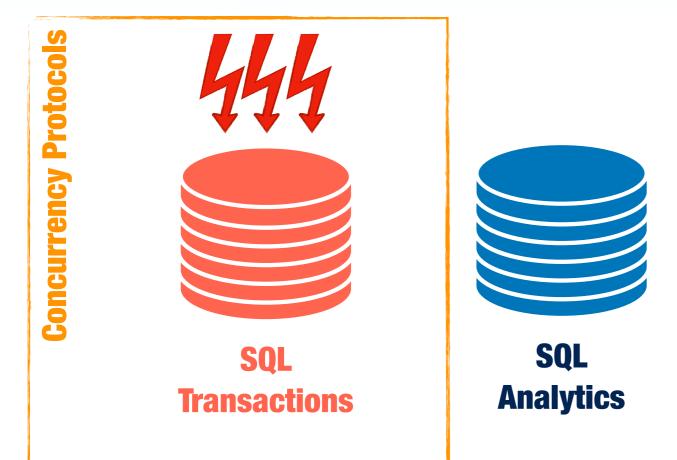
High-dimensional Indexing: (e.g., BE-Tree, BE-topK) [SIGMOD'11, ICDE'12, TODS'13, ICDCS'13, ICDE'14, ICDCS'17, Middleware'17]



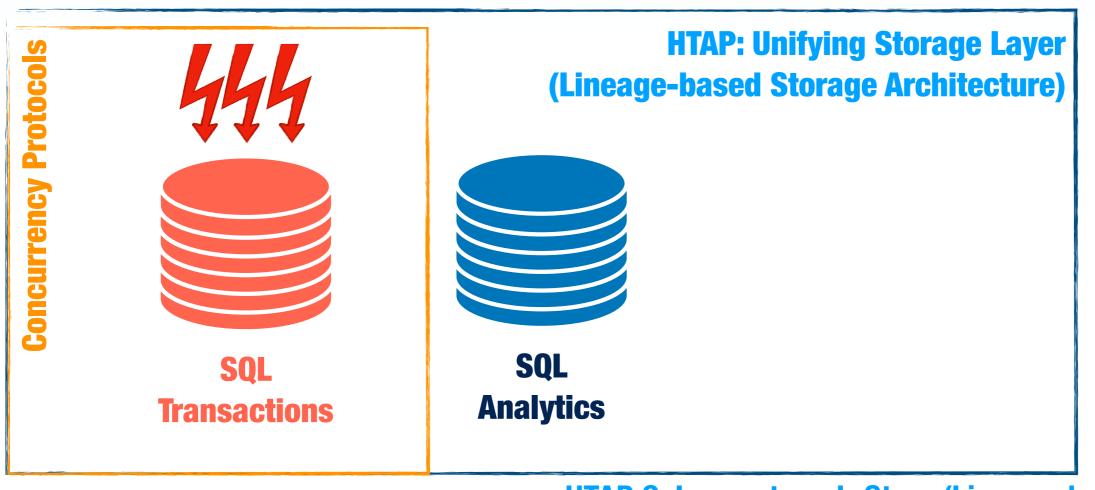
Concurrency Control Protocols: (e.g., 2VCC, QueCC - Best Paper Award) [VLDB'13, VLDB'14, VLDBJ'16, Middleware'16, TDKE'15, SIGMOD'15, ICDE'16, Middleware'18]

QueCC: Queue-Oriented Planning and Execution Architecture



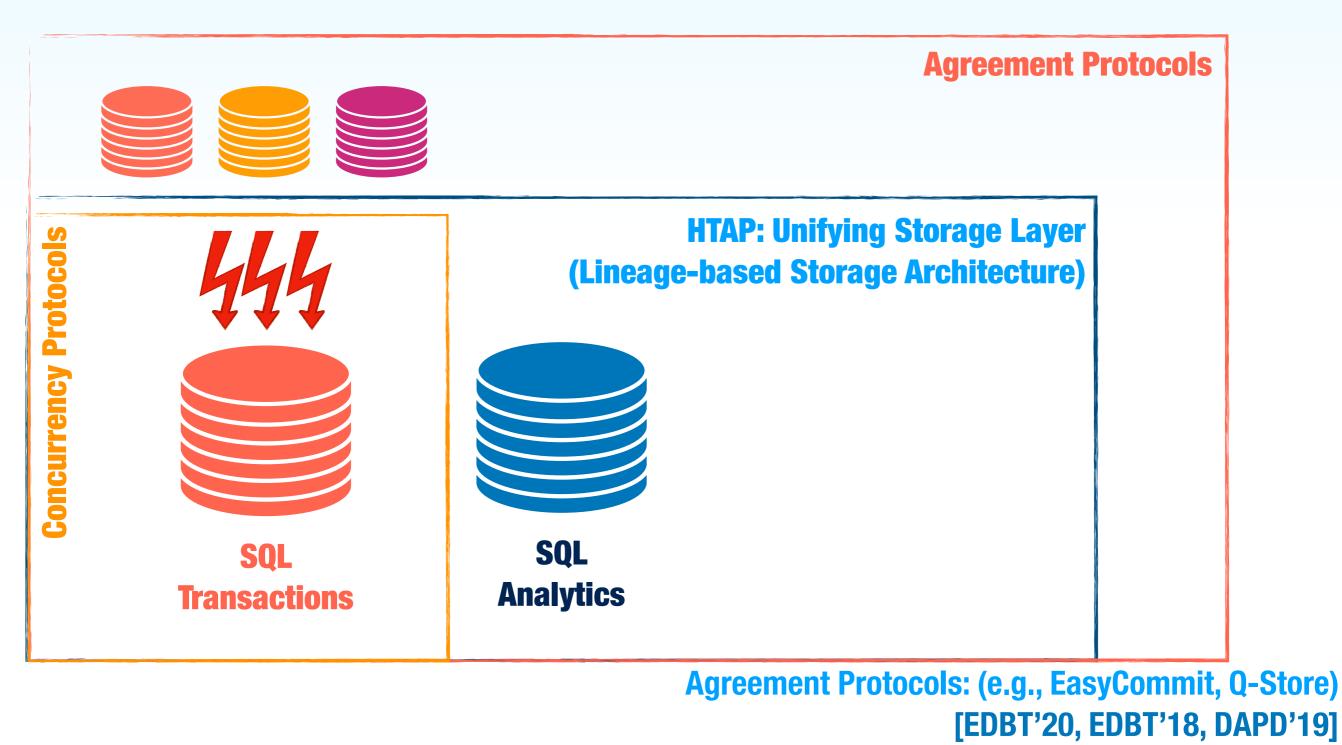


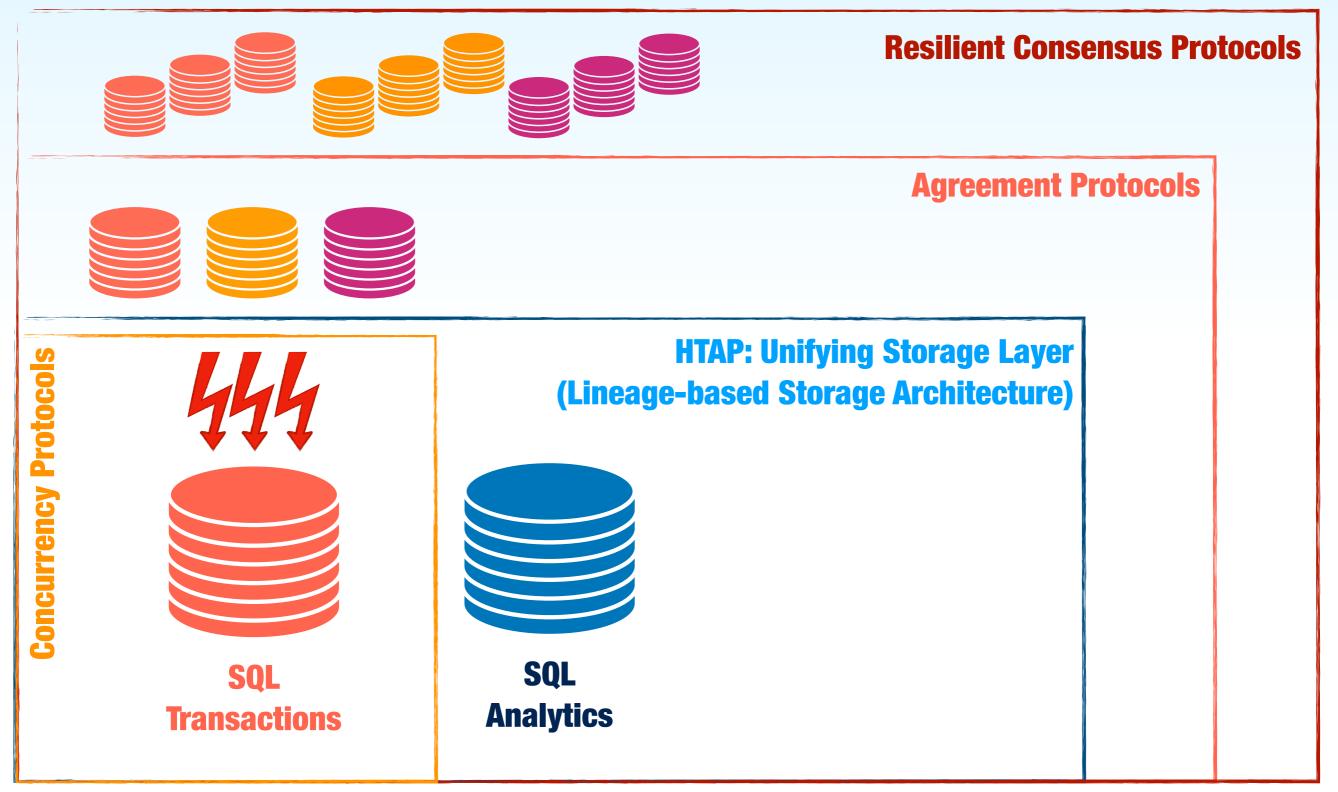
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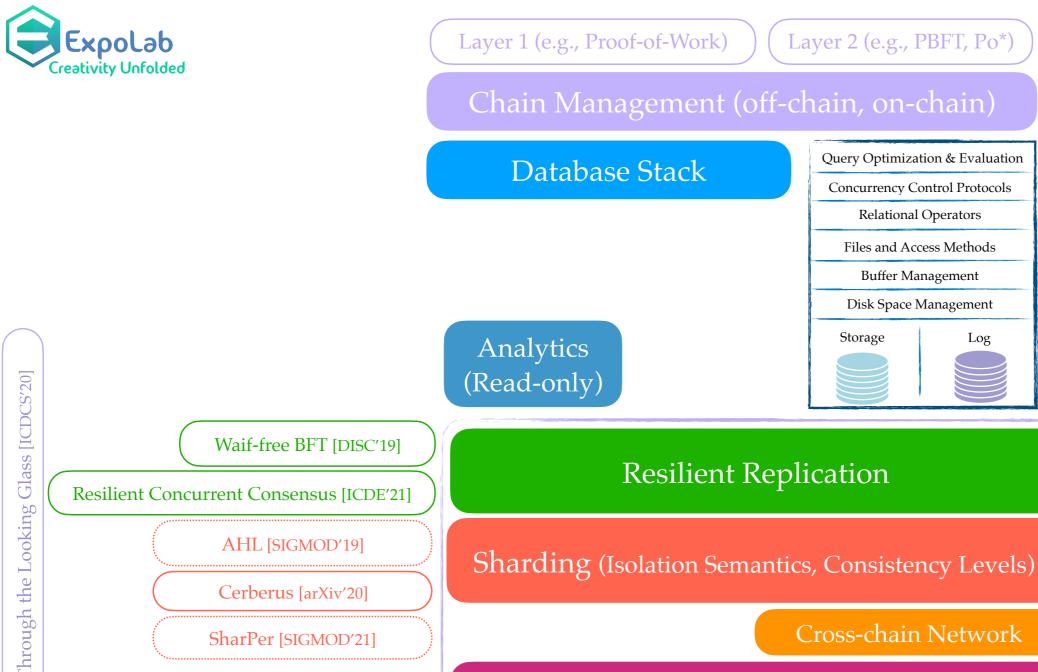
HTAP Column-store: L-Store (Lineage-based Data Store) [VLDB'12, ICDE'14, ICDCS'16, EDBT'18, TKDE'20 (2x) 34 filed US patents]

Graphs on SQL: (e.g., GRFusion) [SIGMOD'18, EDBT'18] 7





Consensus Protocols: (e.g., GeoBFT, PoE, RCC, ByShard, RingBFT, Delayed Replication, CSP, Blockplane) [VLDB'21, ICDE'21, EDTB'21, VLDB'20, ICDCS'20, ICDT'20, DISC'19 (2x), SC'19, ICDE'19, arXiv'19 (8x)]



Global Distribution

Reconfigurable Network

Recovery (View-change)

Identity Management

Permissioned

ResilientDB

curity, Privacy Reloaded

Permissionless

Delayed Replication [ICDT'20]

Proof-of-Execution [EDBT'21]

ByShard [VLDB'21]

RingBFT [arXiv'21]

Atomic Commitment [VLDB'20]

Cross-chain Deals [VLDB'20]

GeoBFT [VLDB'20]

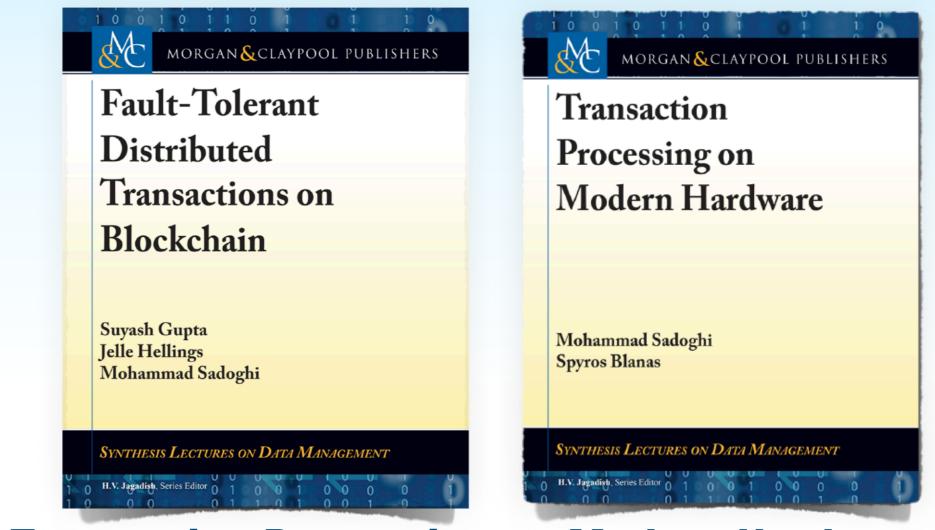
BlockBench [SIGMOD'17]

Cluster Sending Primitive [DISC'19]

Applications: DeFi, Smart Contracts, IoT, Serverless

Blockplane [ICDE'19]





BOOKS Transaction Processing on Modern Hardware.

Synthesis Lectures on Data Management, Morgan & Claypool Publishers 2019

Fault-Tolerant Distributed Transactions on Blockchain.

Synthesis Lectures on Data Management, Morgan & Claypool Publishers 2021









Advancements TV With Ted Danson - CNBC, CityAM, Medium, Yahoo! Finance, Market Insider, CoinDesk, Crypto Media, Davis Enterprise, Times Union, WBOC TV/Radio

Books Transaction Processing on Modern Hardware.

Synthesis Lectures on Data Management, Morgan & Claypool Publishers 2019

Fault-Tolerant Distributed Transactions on Blockchain.

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Quantifiable Resiliency (Graduate Student Experiments)

Aloha Lake, Desolation Wilderness 15 Miles Long 2,500 Feet Elevation Gain (8,700 Feet at Summit)







Tomales Point Trail, Point Reyes National Seashore 9.4 Miles Long 1,579 Feet Elevation Gain

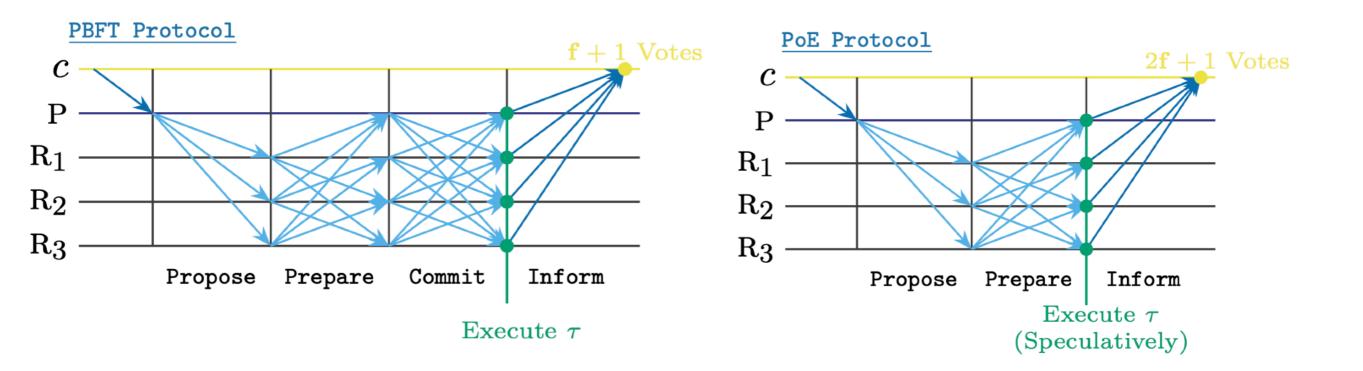




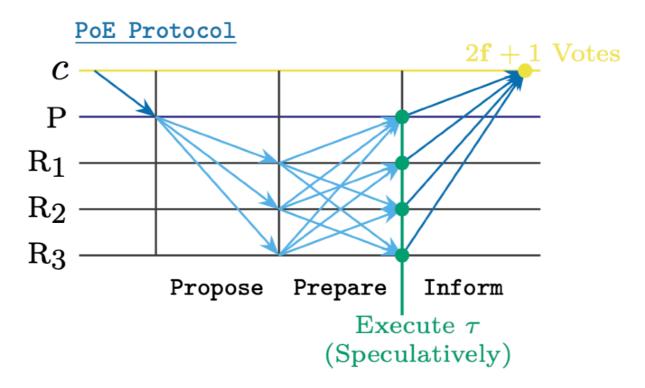
Non-Quantifiable Resiliency

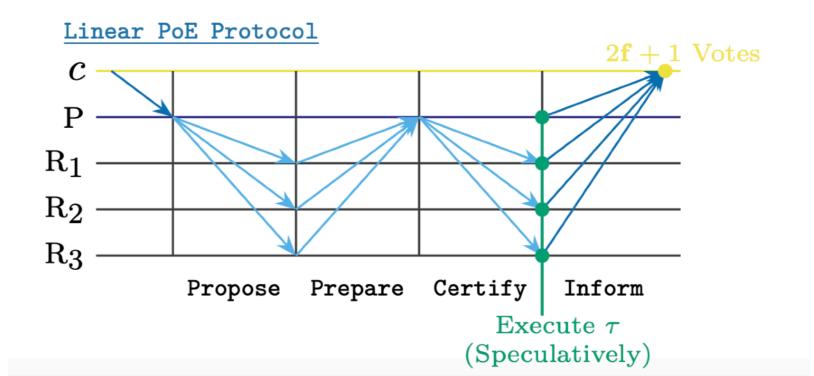
Proof-of-Execution: Reaching Consensus Through Fault-Tolerant Speculation [EDBT'21]

Out-of-Order message processing to reduce replica idleness Speculative Execution with revertible/divergent replicas & eager/irrevertible client commit introducing linear message complexity

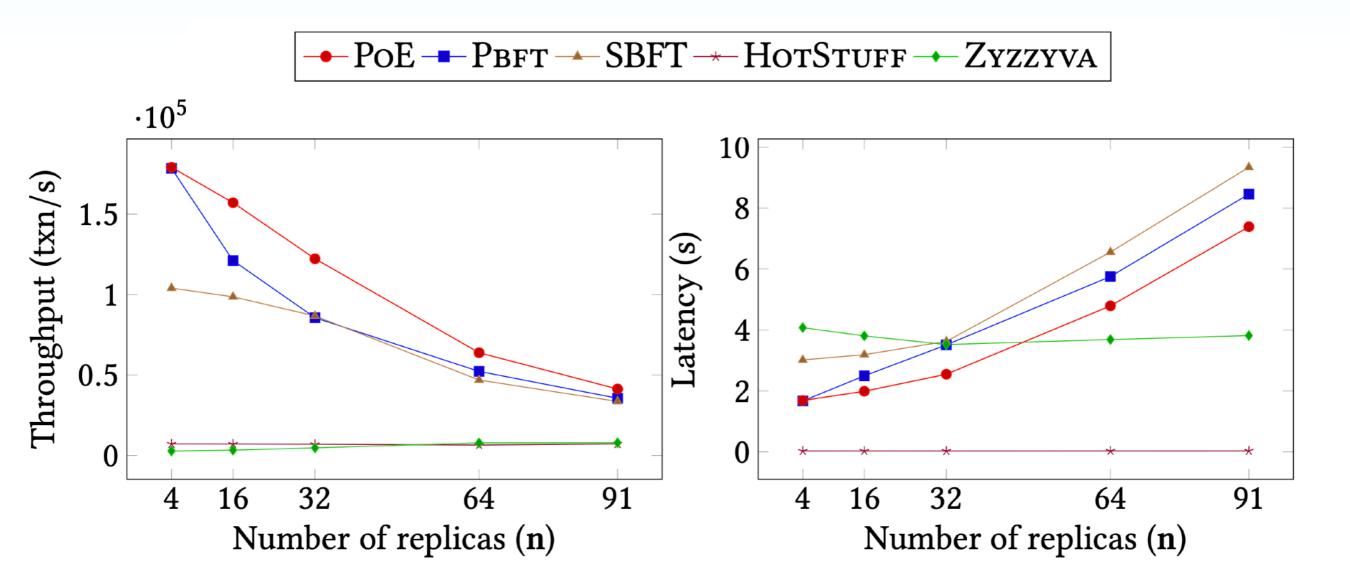


Proof-of-Execution: Reaching Consensus Through Fault-Tolerant Speculation [EDBT'21]





Proof-of-Execution: Reaching Consensus Through Fault-Tolerant Speculation [EDBT'21]



PoE scales beyond <u>91 replicas</u>, in presence of failures, outperforms <u>PBFT up to 43%</u>

RCC: Resilient Concurrent Consensus Paradigm [ICDE'21]

A wait-free meta-protocol...

Designate multiple replicas as primaries!

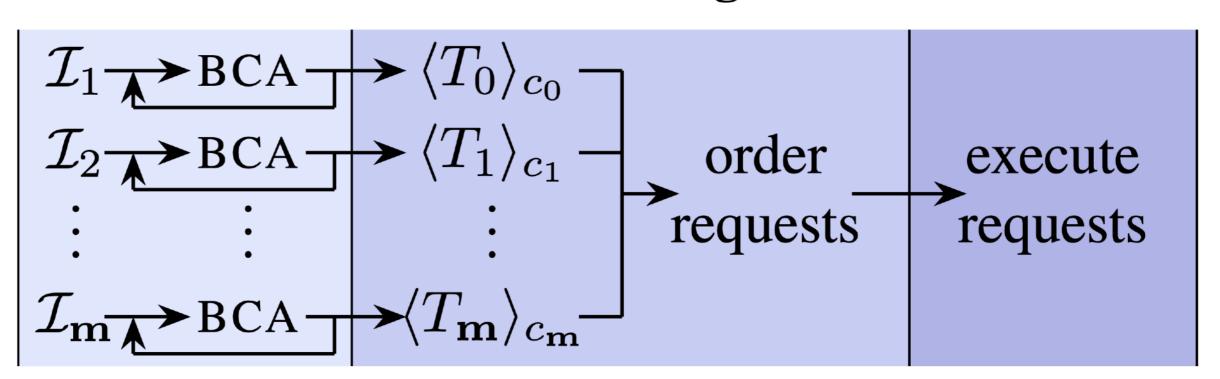
Run multiple parallel consensuses on each replica independently

Concurrent BCA O

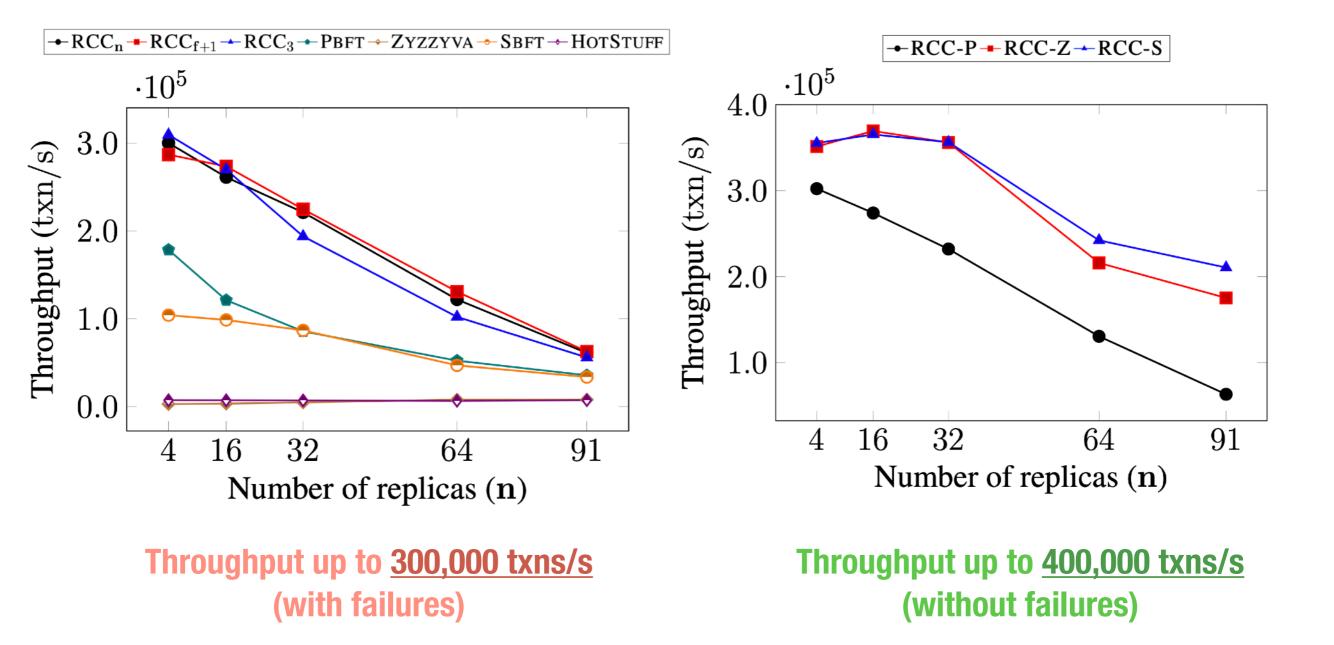
Ordering

Execution

20



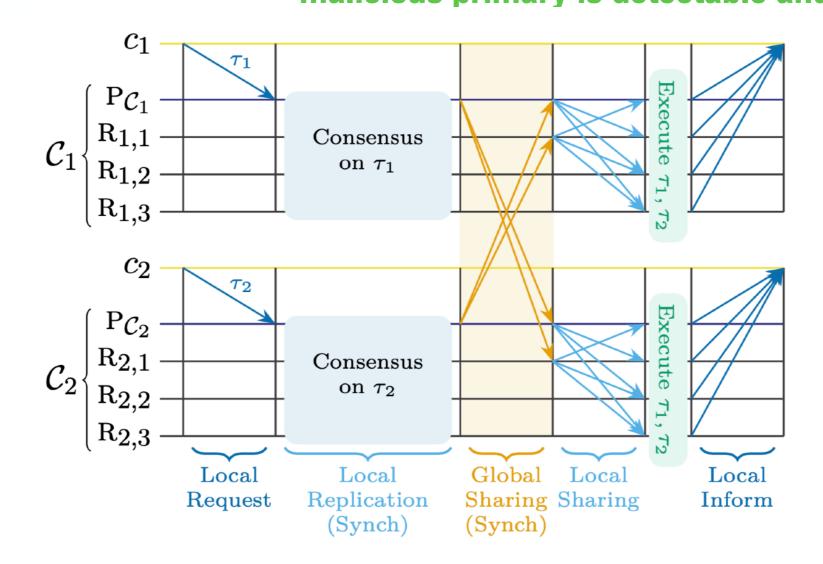
RCC: Resilient Concurrent Consensus Paradigm [ICDE'21]



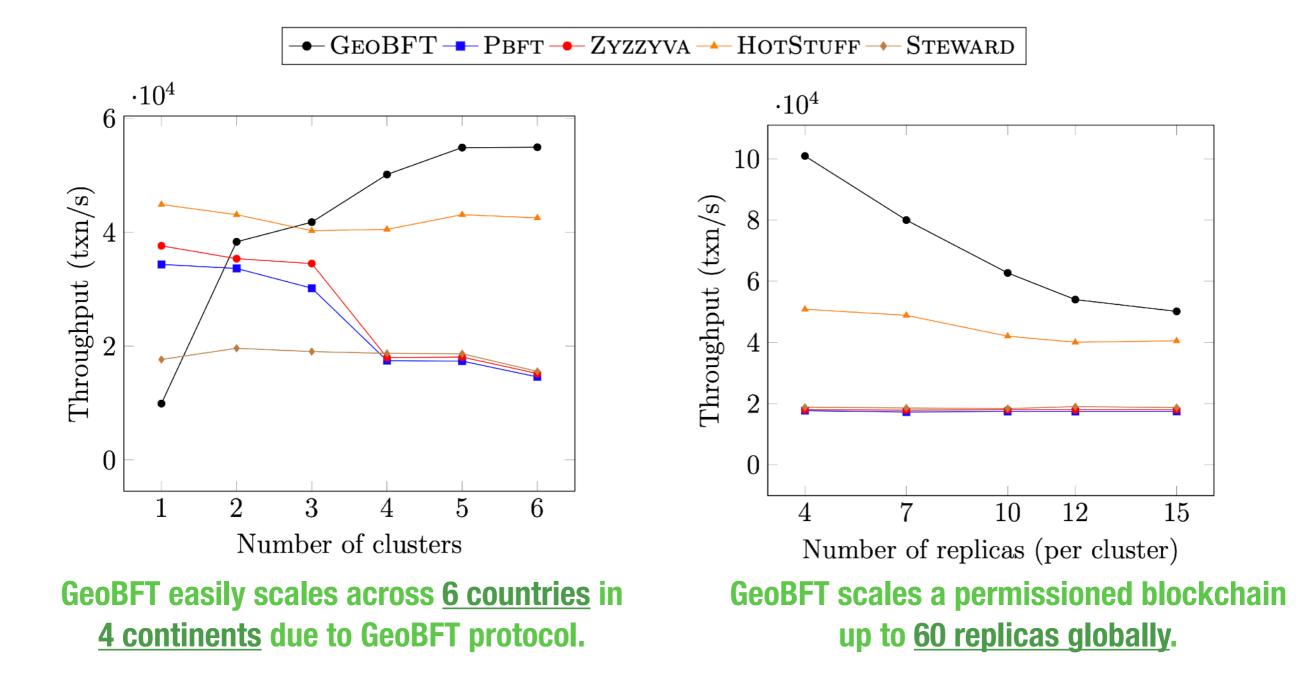
Brief Announcement: Revisiting Consensus Protocols through Wait-free Parallelization. DISC 2019

GeoBFT: Global Scale Resilient Consensus [VLDB'20]

A meta-protocol, locally running any BFT in parallel and independently Global ordering provably requires only linear communication Provably sufficient for primary to send a certificate to at most f+1 replicas, malicious primary is detectable and replaceable



GeoBFT: Global Scale Resilient Consensus [VLDB'20]



The Fault-Tolerant Cluster-Sending Problem [DISC'19]

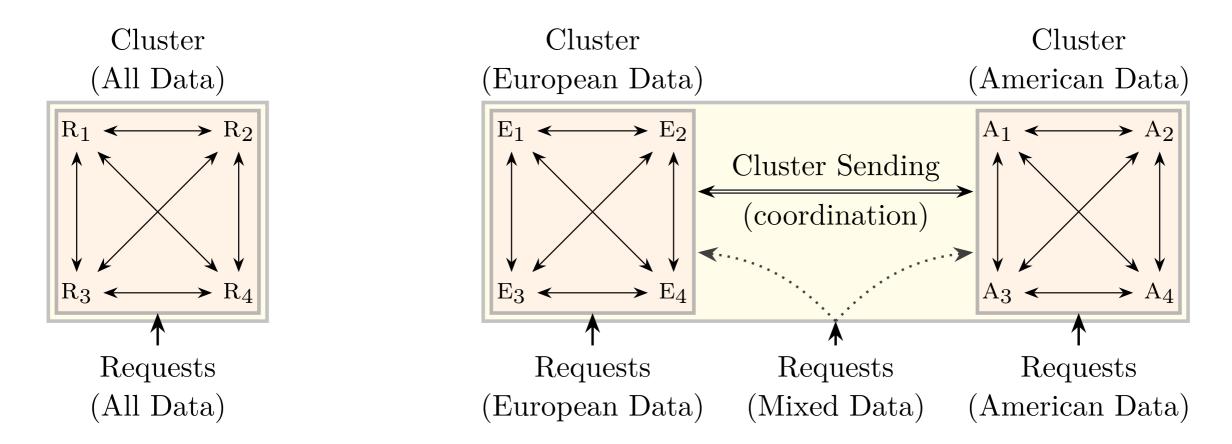
formalizing the problem of sending a message from one Byzantine cluster to

another Byzantine cluster in a reliable manner,

establishing lower bounds on the complexity

of this problem under crash failures and Byzantine failures

(linear in the size of clusters)



The Fault-Tolerant Cluster-Sending Problem [DISC'19]

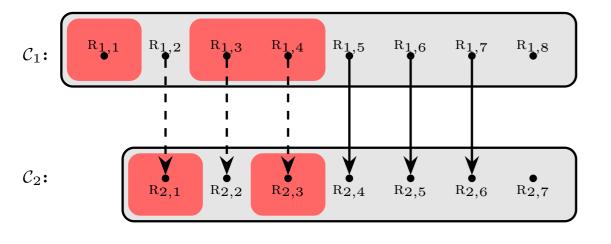
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| | | Protocol | System | Robustness | Messages | Message size |
|---|------------|----------|---------------|---|--|---|
| | non-linear | RB-bcs | Omit | $\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$ | $(\mathbf{f}_{\mathcal{C}_1} + 1) \cdot (\mathbf{f}_{\mathcal{C}_2} + 1)$ | $\mathcal{O}(\ v\)$ |
| | | RB-brs | Byzantine, RS | $\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$ | $(2\mathbf{f}_{\mathcal{C}_1}+1)\cdot(\mathbf{f}_{\mathcal{C}_2}+1)$ | $\mathcal{O}(\ v\)$ |
| | | RB-bcs | Byzantine, RS | $\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$ | $(\mathbf{f}_{\mathcal{C}_1} + 1) \cdot (\mathbf{f}_{\mathcal{C}_2} + 1)$ | $\mathcal{O}(\ v\ + \mathbf{f}_{\mathcal{C}_1})$ |
| | | RB-bcs | Byzantine, CS | $\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$ | $(\mathbf{f}_{\mathcal{C}_1}+1)\cdot(\mathbf{f}_{\mathcal{C}_2}+1)$ | $\mathcal{O}(\ v\)$ |
| - | linear | PBS-bcs | Omit | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | $\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})) \text{ (optimal)}$ | $\mathcal{O}(\ v\)$ |
| | | PBS-brs | Byzantine, RS | $\mathbf{n}_{\mathcal{C}_1} > 4\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 4\mathbf{f}_{\mathcal{C}_2}$ | $\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})) 	ext{ (optimal)}$ | $\mathcal{O}(\ v\)$ |
| | | PBS-bcs | Byzantine, RS | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | $\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2}))$ | $\mathcal{O}(\ v\ + \mathbf{f}_{\mathcal{C}_1})$ |
| | | PBS-bcs | Byzantine, CS | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | $\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})) 	ext{ (optimal)}$ | $\mathcal{O}(\ v\)$ |

Brief Announcement: The Fault-Tolerant Cluster-sending Problem. DISC 2019

Byzantine Cluster-Sending in Expected Constant Communication [arXiv'21]

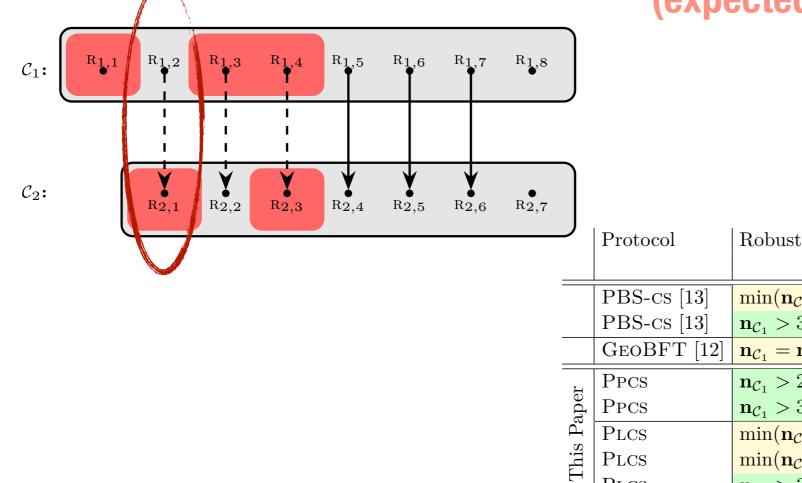
formalizing the problem of probabilistically sending a message from one

Byzantine cluster to another Byzantine cluster in a reliable manner,

establishing lower bounds on the complexity

of this problem under crash failures and Byzantine failures

(expected constant message complexity)



| | Protocol | Robustness | Message Steps | | | U. |
|-----------|---------------|---|---|--|--------------|----|
| | | | (expected) | (worst) | | |
| | PBS-cs [13] | $\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$ | $\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$ | | | × |
| | PBS-cs [13] | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | $\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$ | | | × |
| | GeoBFT $[12]$ | $\mathbf{n}_{\mathcal{C}_1} = \mathbf{n}_{\mathcal{C}_2} > 3\max(\mathbf{f}_{\mathcal{C}_1}, \mathbf{f}_{\mathcal{C}_2})$ | $\mathbf{f}_{\mathcal{C}_2} + 1^{\ddagger}$ | $\Omega(\mathbf{f}_{\mathcal{C}_1}\mathbf{n}_{\mathcal{C}_2})$ | × | ✓ |
| nus raper | PPCS | $\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 2\mathbf{f}_{\mathcal{C}_2}$ | 4 | $(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$ | × | ✓ |
| | Ppcs | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | $2\frac{1}{4}$ | $(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$ | × | ✓ |
| | PLCS | $\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$ | 4 | $\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$ | < | ~ |
| | Plcs | $\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > 2(\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2})$ | $2\frac{1}{4}$ | $\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$ | \checkmark | ~ |
| | PLCS | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | 3 | $\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$ | < | ✓ |

Byzantine Cluster-Sending in Expected Constant Communication [arXiv'21]

formalizing the problem of probabilistically sending a message from one

Byzantine cluster to another Byzantine cluster in a reliable manner,

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(expected constant message complexity)

| \mathcal{C}_1 : | $\begin{array}{c c} R_{1,1} \\ \bullet \end{array} \begin{array}{c} R_{1,2} \\ \bullet \end{array} \begin{array}{c} R_{1,3} \\ \bullet \end{array} \begin{array}{c} R_{1,4} \\ \bullet \end{array}$ | R _{1,5} R _{1,6} | R _{1,7} R _{1,8} | 8 | | | | | | |
|-------------------|---|-----------------------------------|-----------------------------------|-------|-------------|---|---|---|----------|----|
| | | | | | | | | | | |
| \mathcal{C}_2 : | R _{2,1} R _{2,2} R _{2,3} | R _{2,4} R _{2,5} | R _{2,6} R _{2,7} | 7 | Protocol | Robustness | Mes (expected) | ssage Steps (worst) | 0. | U. |
| | | | | | PBS-cs [13] | $\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$ | \mathbf{f}_{C_1} | $+\mathbf{f}_{C_2}+1$ | ~ | × |
| | | | | | | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | - | $\mathbf{x}(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$ | ~ | × |
| | | | | | GEOBFT [12] | $\mathbf{n}_{\mathcal{C}_1} = \mathbf{n}_{\mathcal{C}_2} > 3\max(\mathbf{f}_{\mathcal{C}_1}, \mathbf{f}_{\mathcal{C}_2})$ | $\mathbf{f}_{\mathcal{C}_2} + 1^{\ddagger}$ | $\Omega(\mathbf{f}_{\mathcal{C}_1}\mathbf{n}_{\mathcal{C}_2})$ | × | ✓ |
| | | | | | PPCS | $\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 2\mathbf{f}_{\mathcal{C}_2}$ | 4 | $(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$ | X | ✓ |
| | | | | Paper | PPCS | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | $2\frac{1}{4}$ | $(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$ | × | ✓ |
| | | | | s D | | $\min(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$ | 4 | $\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$ | ✓ | |
| | | | | This | PLCS | $\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > 2(\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2})$ | $2\frac{1}{4}$ | $\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$ | ✓ | |
| | | | | L 1 | PLCS | $\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$ | 3 | $\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$ | ✓ | ✓ |

RingBFT: Resilient Consensus Over Sharded Ring Topology [arXiv'21]

A meta-protocol adhering to the ring order, and follow the principle of

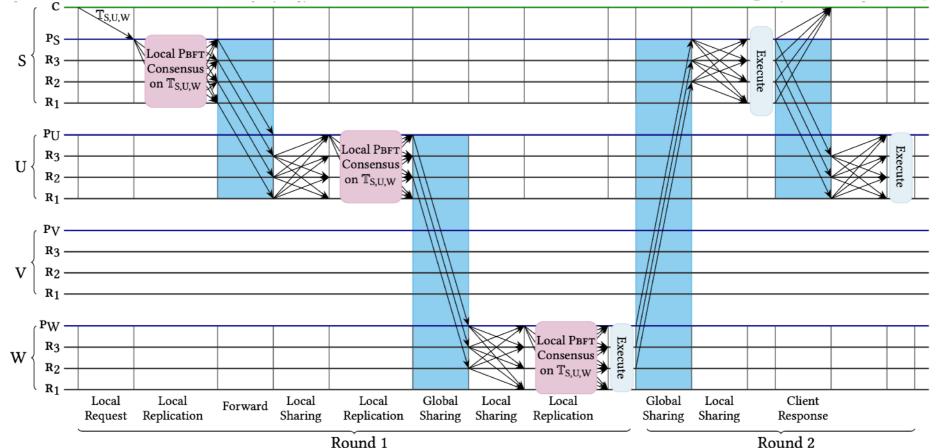
process, forward, and re-transmit

Guarantees consensus for each cross-shard transaction in

at most two rotations around the ring

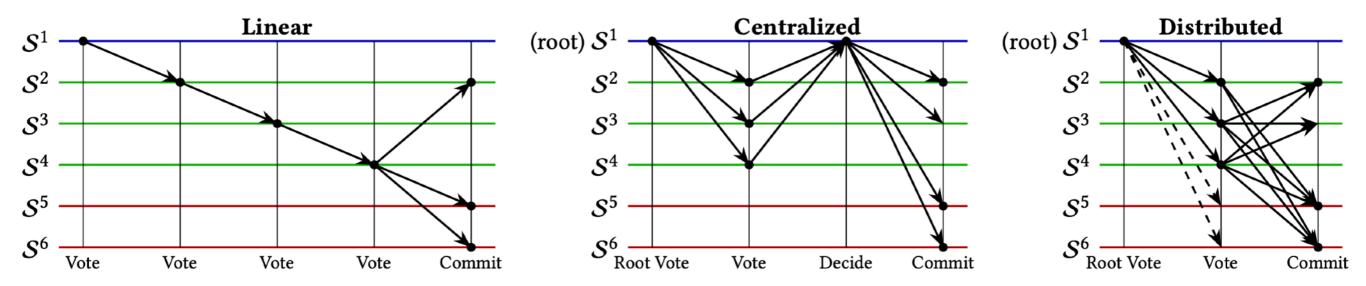
Sustaining over 1,200,000 transactions per second when deployed globally spanning ten

countries, fifteen regions, nearly 500 replicas.



ByShard: Sharding in a Byzantine Environment [VLDB'21]

Processing multi-shard transaction via the orchestrate-execute model Processing is broken down into three types of shard-steps: vote, commit, and abort Each shard-step is performed via one consensus step Steps are communicated via cluster-sending



Coordination-Free Byzantine Replication With Minimal Communication Costs [ICDT'20]

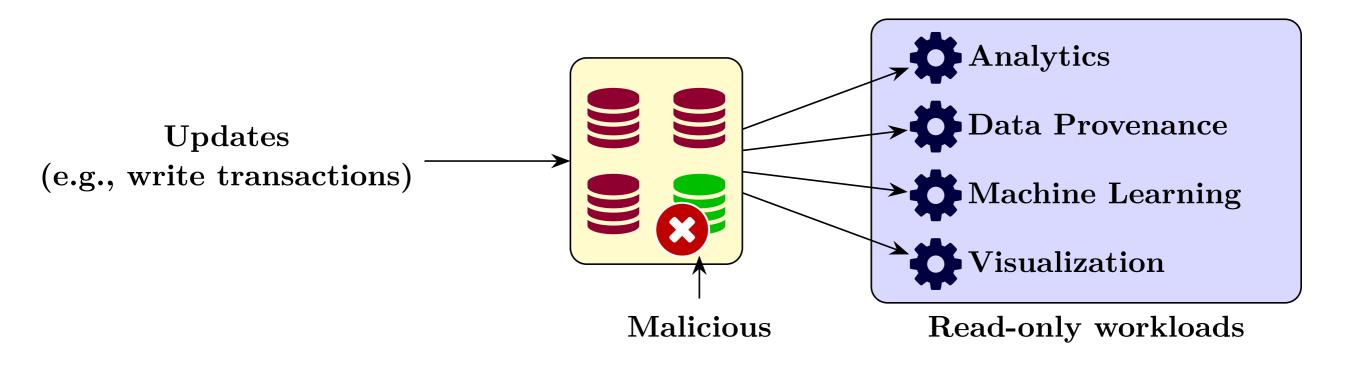
formalizing the Byzantine learner problem to support efficient

analytics for blockchain applications

introducing the delayed-replication algorithm,

utilizing information dispersal techniques,

giving rise to a coordination-free, push-based, minimal communication protocol



Coordination-Free Byzantine Replication With Minimal Communication Costs [ICDT'20]

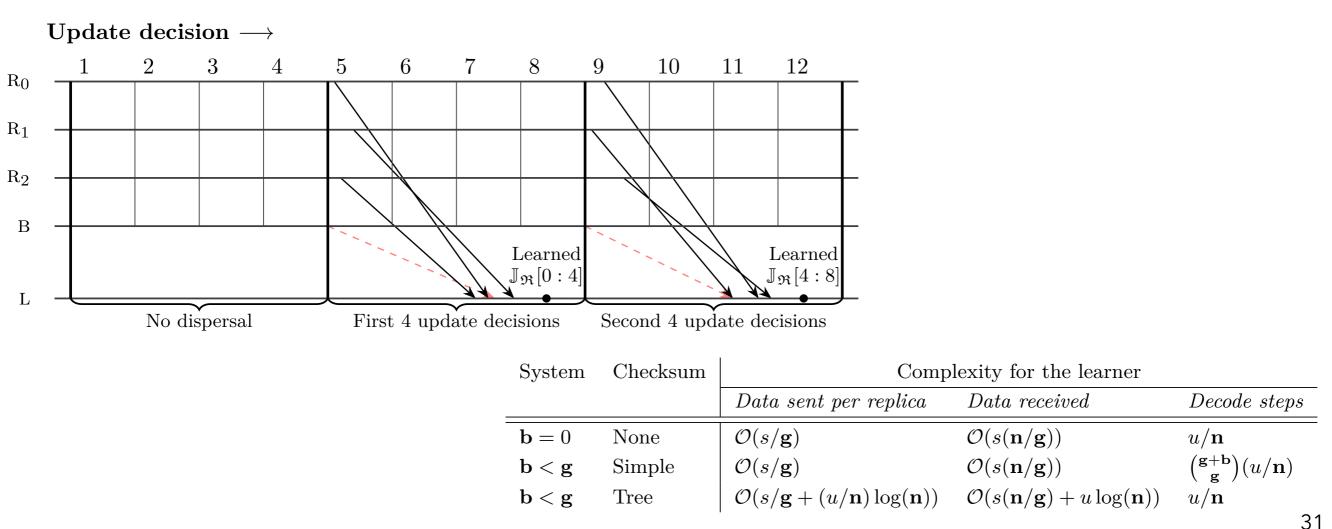
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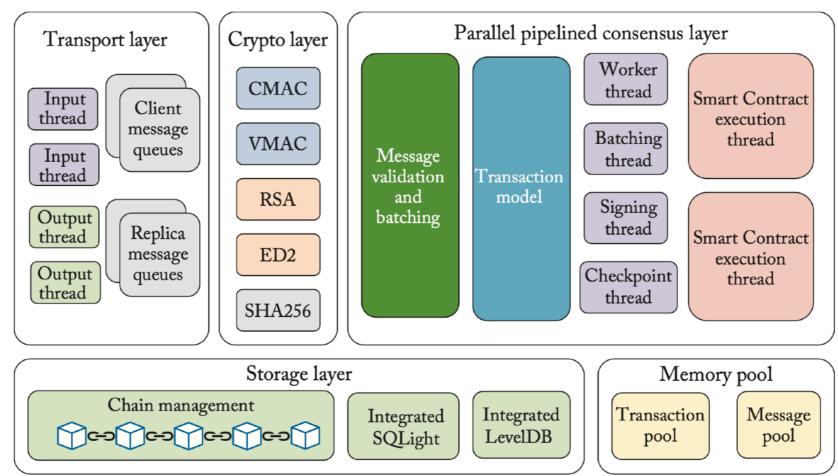
utilizing information dispersal techniques,

giving rise to a coordination-free, push-based, minimal communication protocol



Permissioned Blockchain Through the Looking Glass: Architectural and Implementation Lessons Learned [ICDCS'20]

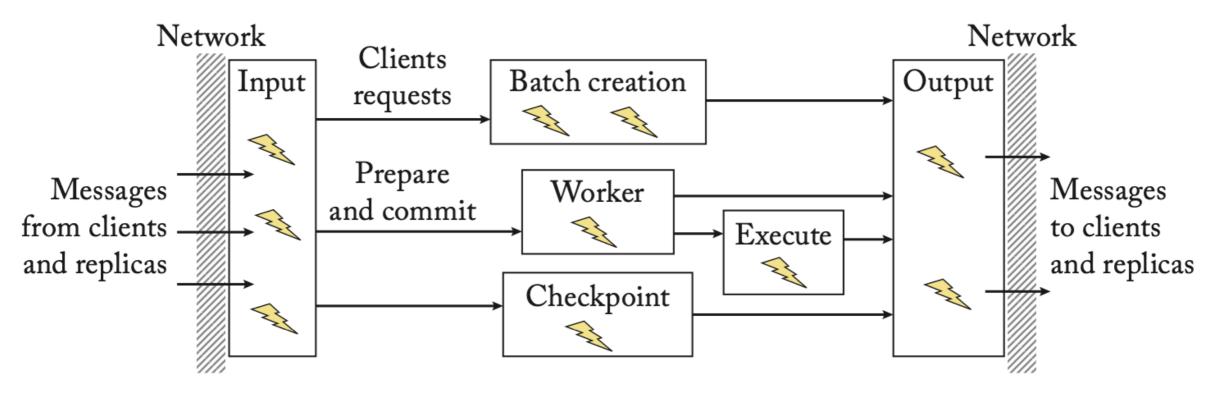
Single-threaded Monolithic Design Out-of-ordering Consensus Communication De-coupled Ordering and Execution Off-Chain Memory Management Expensive Cryptographic Practices (DS vs. MAC) Smart Contracts Code Generation (Pre-compilation)





Permissioned Blockchain Through the Looking Glass: Architectural and Implementation Lessons Learned [ICDCS'20]

Single-threaded Monolithic Design Out-of-ordering Consensus Communication De-coupled Ordering and Execution Off-Chain Memory Management Expensive Cryptographic Practices (DS vs. MAC) Smart Contracts Code Generation (Pre-compilation)

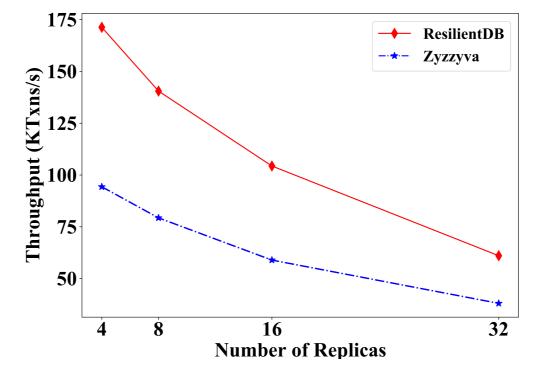


Multi-Threaded Deep Pipeline

Security, Privacy Reloaded

Permissioned Blockchain Through the Looking Glass: Architectural and Implementation Lessons Learned [ICDCS'20]

Single-threaded Monolithic Design Out-of-ordering Consensus Communication De-coupled Ordering and Execution Off-Chain Memory Management Expensive Cryptographic Practices (DS vs. MAC) Smart Contracts Code Generation (Pre-compilation)





Can a well-crafted system based on a classical BFT protocol outperform a modern protocol?

Revisit Resiliency (Graduate Student Experiment Continues)

Mount Tallac, Lake Tahoe 12.1 Miles Long 3,931 Feet Elevation Gain (9,738 Feet at Summit)





Fostering Resiliency

(Offering Stress Management and Well-Being Courses at UC Davis)



THE CALIFORNIA AGGIE Seminar spotlight: "Becoming an Extraordinary Human"

Spring 2020

Tamarkoz®.

ECS 298 (CRN 66553):

Days: Wednesdays

INSTRUCTORS:

Mohammad Sadoghi, Ph.D. Nasim Bahadorani, DrPH.

Time: 7:00 pm - 8:00 pm

Graduate Survival Kit

Learn the foundation & working

knowledge of stress reduction based on a unique heart-centered

meditation practice referred to as

The M.T.O. Tamarkoz® method is

the art of self-knowledge through

concentration and meditation.

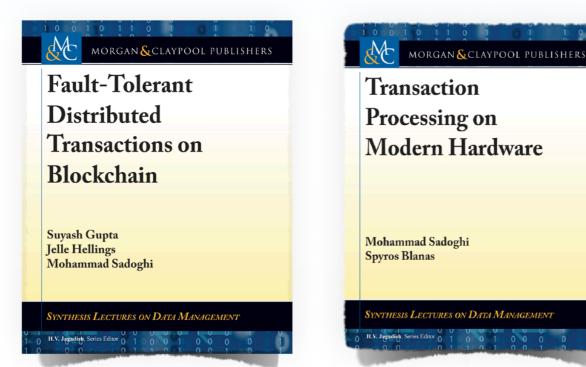
The California Aggie, April 6, 2020

BE BALANCED





THANK YOU



FOR COMPLETE REFERENCES



Google